

Braid monodromy computations using certified path tracking

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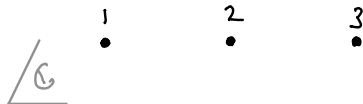
Warmup: braid group

Definitions

$$OC_n = \{(x_1, \dots, x_n) \in \mathbb{C}^n : \forall i \neq j, x_i \neq x_j\}.$$

$$C_n = \{\text{subsets of size } n \text{ in } \mathbb{C}\}.$$

$$(x_1, \dots, x_n) \in OC_n \mapsto \{x_1, \dots, x_n\} \in C_n.$$



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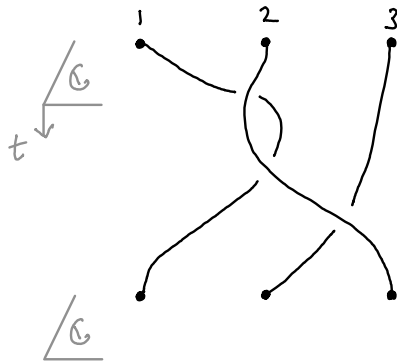
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A braid is a homotopy class of a path

$\beta : [0, 1] \rightarrow C_n$ such that

$\beta(0) = \beta(1) = \{1, \dots, n\}$,



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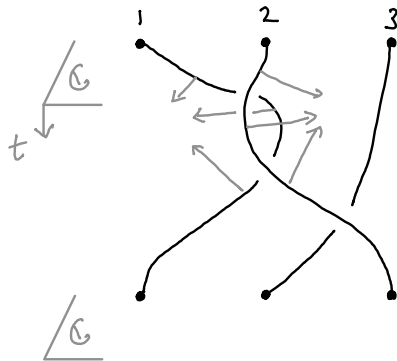
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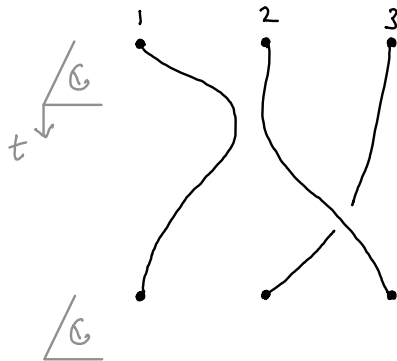
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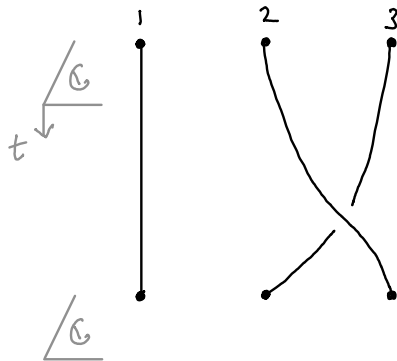
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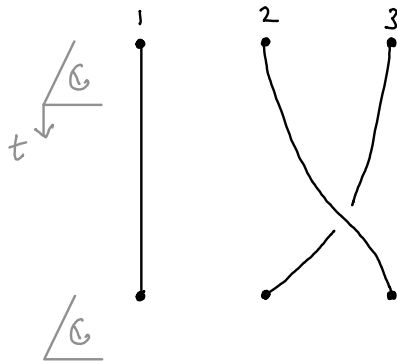
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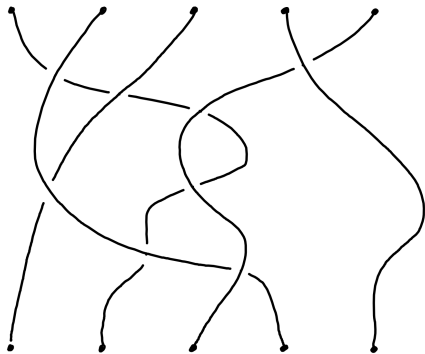
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Remark

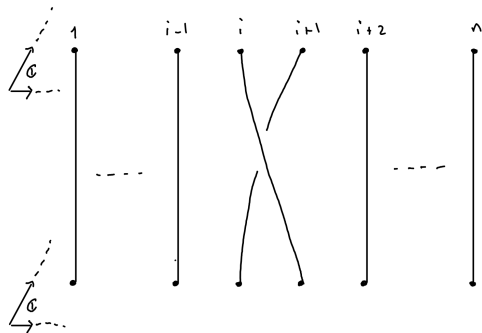
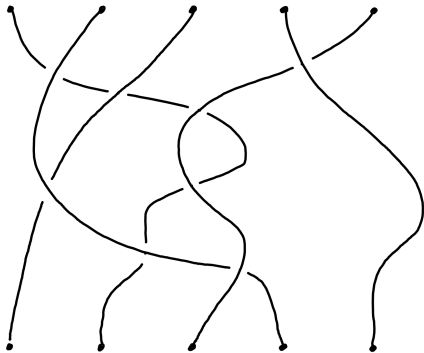
A path $\zeta = (\zeta_1, \dots, \zeta_n) : [0, 1] \rightarrow OC_n$ induces a braid. **If $\zeta' : [0, 1] \rightarrow OC_n$ is homotopic to ζ , they have the same associated braid**



Examples and Artin's theorem

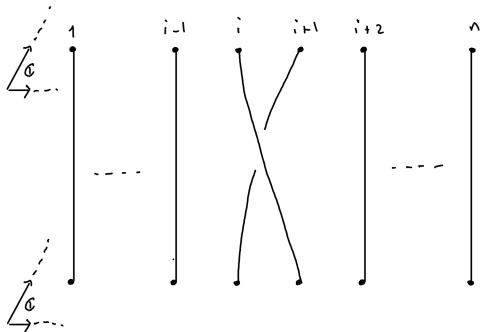
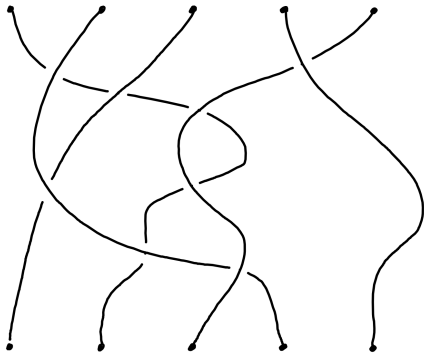


Examples and Artin's theorem



Standard generator σ_i

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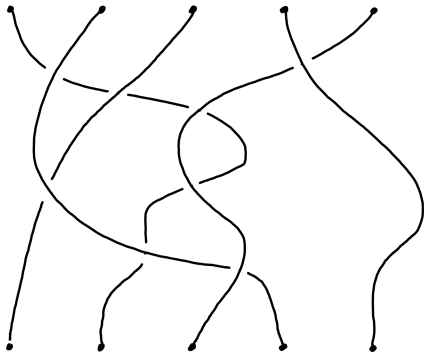


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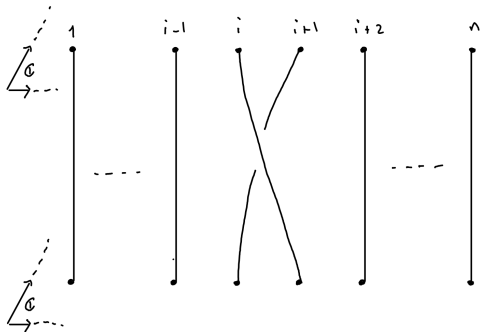
Theorem [Artin, 1947]

The σ_i 's generate B_n (+ explicit relations).

Examples and Artin's theorem



$$\sigma_4 \sigma_1^{-1} \sigma_2^{-1} \sigma_3^{-1} \sigma_3 \sigma_1 \sigma_2 \sigma_3^{-1}$$

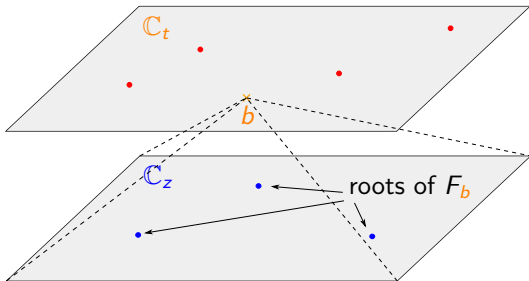


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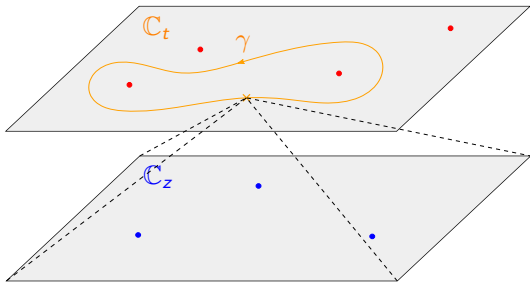
Another example



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- define $F_t(z) = g(t, z)$.
- Let $b \in \mathbb{C} \setminus \Sigma$ be a base point,

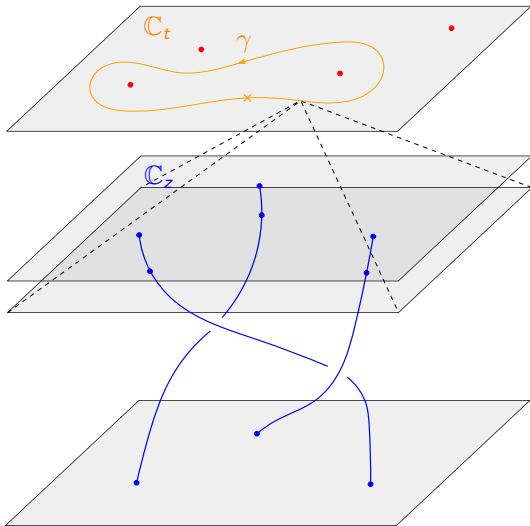
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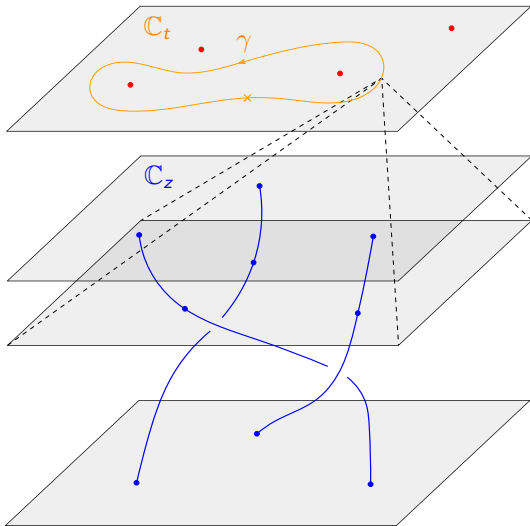
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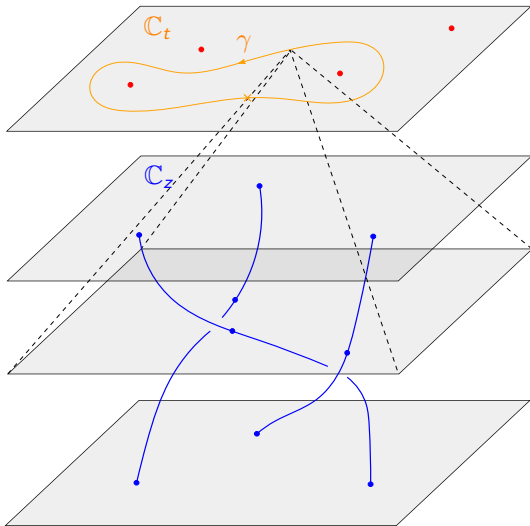
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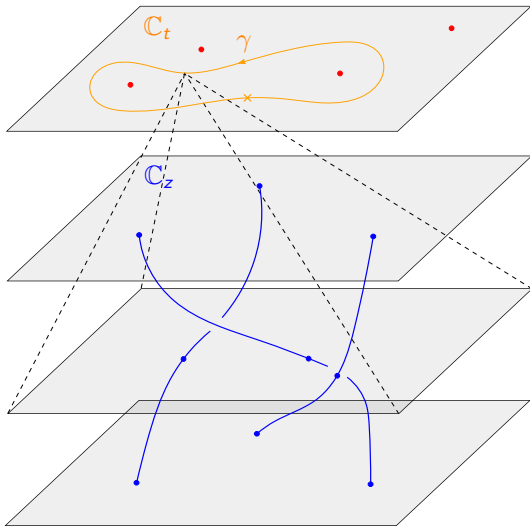
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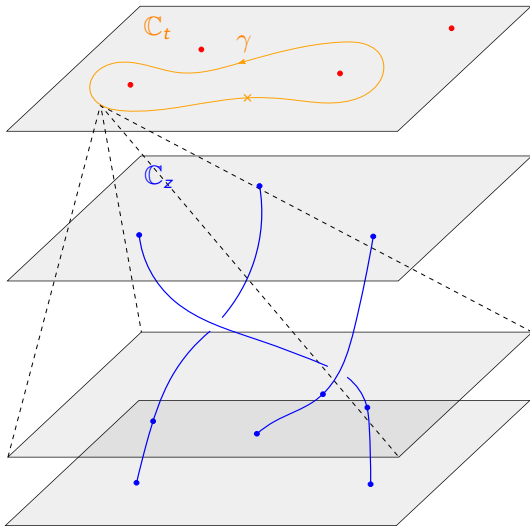
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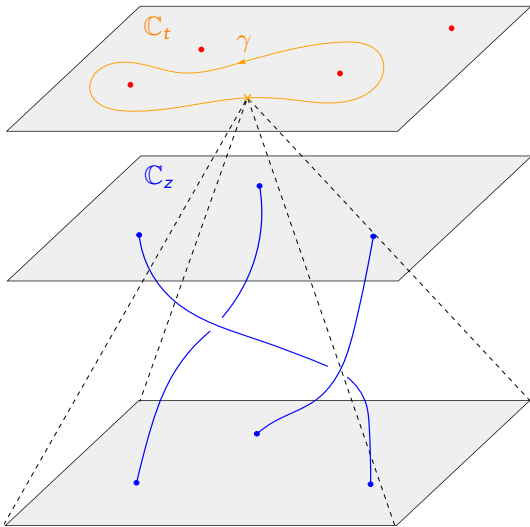
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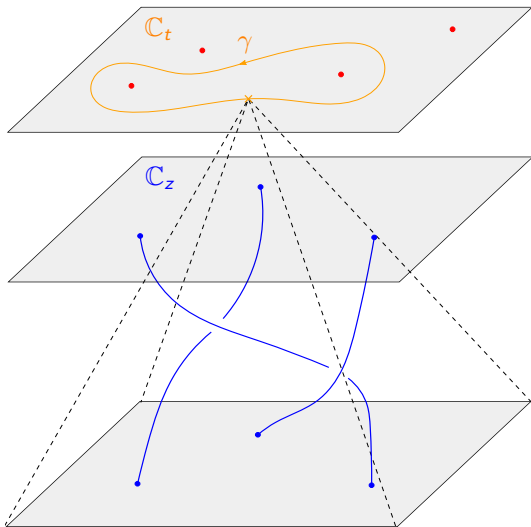
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Algorithmic goal

Input: g, γ

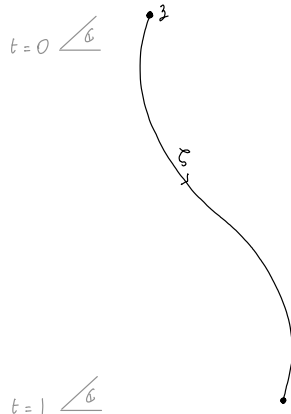
Output: the associated braid in terms of Artin's generators

Main tool

Certified homotopy continuation

Input: $H : [0, 1] \times \mathbb{C}^r \rightarrow \mathbb{C}^r$ and $z \in \mathbb{C}^r$ such that $H(0, z) = 0$.

There exists $\zeta : [0, 1] \rightarrow \mathbb{C}^r$ such that $H(t, \zeta(t)) = 0$ and $\zeta(0) = z$. Assume it is unique.



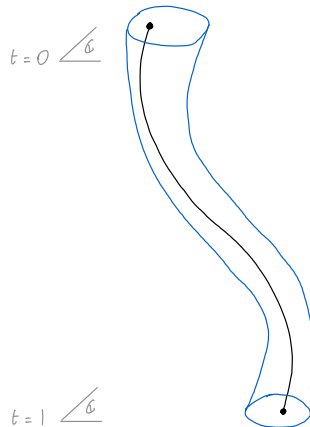
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Output: A tubular neighborhood isolating ζ .



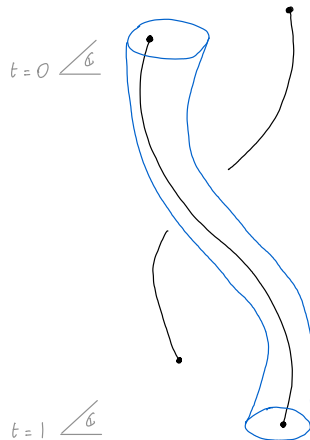
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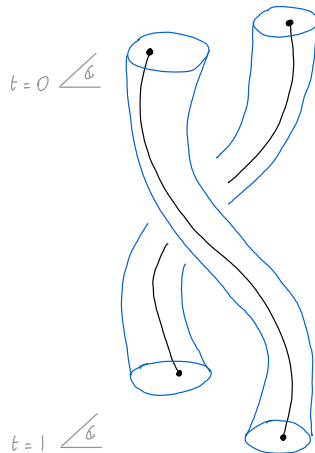
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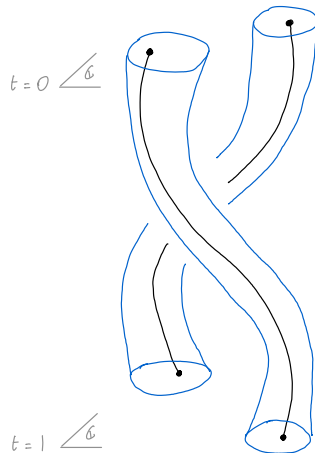
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Application

Take $g \in \mathbb{C}[t, z]$ from last slide and $\gamma : [0, 1] \rightarrow \mathbb{C} \setminus \Sigma$ ($n = \deg_z(g)$). Apply certified homotopy continuation to $H(t, z) = g(\gamma(t), z)$.

Goal: use Algpah [G. and Lairez, 2024] for this step



Certified homotopy continuation

- Kearfott, R. B., & Xing, Z. (1994). An Interval Step Control for Continuation Methods.
- van der Hoeven, J. (2015). *Reliable homotopy continuation*.
- Xu, J., Burr, M., & Yap, C. (2018). An Approach for Certifying Homotopy Continuation Paths: Univariate Case.
- Duff, T., & Lee, K. (2024). Certified homotopy tracking using the Krawczyk method.

Braid computations

- Rodriguez, J. I., & Wang, B. (2017). [Numerical computation of braid groups](#).
- Marco-Buzunariz, M. Á., & Rodríguez, M. (2016). [SIROCCO: a library for certified polynomial root continuation](#).

Braid algorithm

We now assume $\zeta = (\zeta_1, \dots, \zeta_n) : [0, 1] \rightarrow OC_n$.

Goal

Input : n disjoint tubular neighborhoods around ζ_1, \dots, ζ_n

Output : A decomposition in standard generators of the braid induced by ζ_1, \dots, ζ_n

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Interface

We assume a function $\text{sep}(i, j, t)$ that returns $t' \in (t, 1]$ and a symbol in $\star \in \{\rightarrow, \leftarrow, \rightarrow, \leftarrow\}$, such that for all $s \in [t, t']$,

- $\text{Re}(\zeta_i(s)) < \text{Re}(\zeta_j(s))$ if $\star = \rightarrow$,
- $\text{Re}(\zeta_i(s)) > \text{Re}(\zeta_j(s))$ if $\star = \leftarrow$,
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Easy to implement in practice thanks to interval arithmetic !

Cells

Recall: $OC_n = \{(x_1, \dots, x_n) \in \mathbb{C}^n : \forall i \neq j, x_i \neq x_j\}$.

Definition

A cell is a pair $c = (R, I)$ of relations on $\{1, \dots, n\}$.

We associate to it a topological space $|c| \subseteq OC_n$

which points are $(x_1, \dots, x_n) \in OC_n$ such that

- for all $(i, j) \in R$, $\operatorname{Re}(x_i) < \operatorname{Re}(x_j)$,
- for all $(i, j) \in I$, $\operatorname{Im}(x_i) < \operatorname{Im}(x_j)$,

Notation

- $i \xrightarrow{\text{blue}} j \iff (i, j) \in R$
- $i \xrightarrow{\text{red}} j \iff (i, j) \in I$

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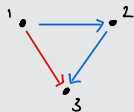
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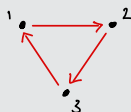
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Examples

$c = (\emptyset, \emptyset)$: $|c| = OC_n$,



$(1, 2, 3 + i) \notin |c|$



$|c| = \emptyset$

Properties of cells

Empty cells

A cell is empty if and only if there is a cycle in R or in I .

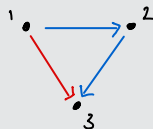
Convex cells

A (non-empty) cell is convex if and only if for all $i, j \in \{1, \dots, n\}$, either $i \rightarrow^* j$ or $j \rightarrow^* i$ or $i \rightarrow^* j$ or $j \rightarrow^* i$. We call this graph property “monochromatic semi-connectedness” (m.s.c. for short).

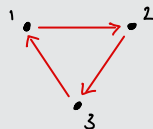
Intersection of cells

Given $c = (R, I)$ and $c' = (R', I')$ two cells, the space associated to $(R \cup R', I \cup I')$ is $|c| \cap |c'|$.

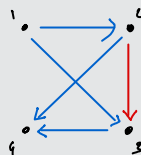
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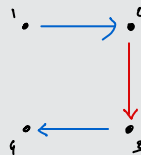
$|c| \neq \emptyset$



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$|c|$ convex

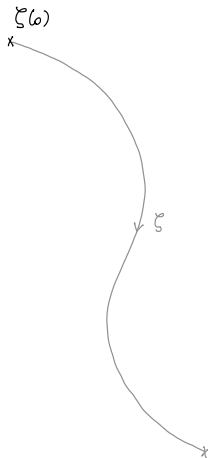


$|c|$ not convex

Step 1: compute a sequence of cells

Path to cells

Input: ζ (represented by tubular neighborhoods)



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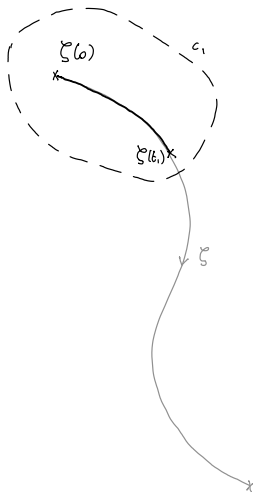
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Idea:

- Start with an initial m.s.c cell c containing $\zeta(0)$.
- Associate to each edge a time of validity.
- When a relation expires, update it using sep and repair m.s.c.
- Repeat



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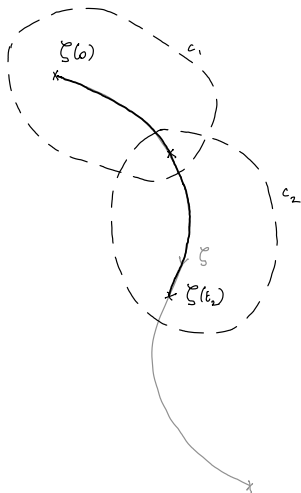
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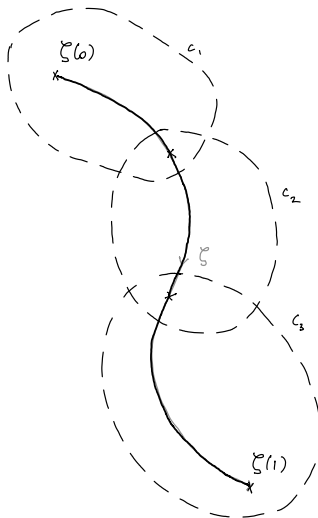
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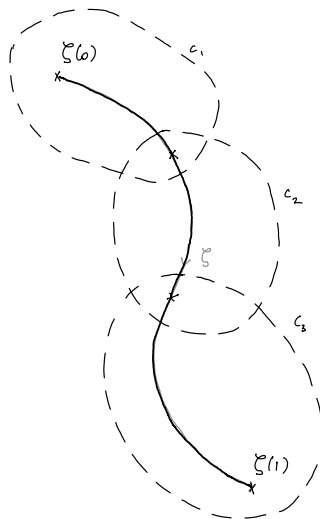
Step 2: linearize ζ

Definition

Let $\rho, \iota \in \mathfrak{S}_n$. We define

$$\omega_{\rho, \iota} = (\rho(1) + \mathbf{i}\iota(1), \dots, \rho(n) + \mathbf{i}\iota(n)) \in OC_n$$

$$\omega_{(12)(34), (143)} = \begin{array}{cccc} \cdot & \bullet_1 & \cdot & \cdot \\ \cdot & \cdot & \bullet_4 & \cdot \\ \bullet_2 & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \bullet_3 \end{array}$$



Step 2: linearize ζ

Definition

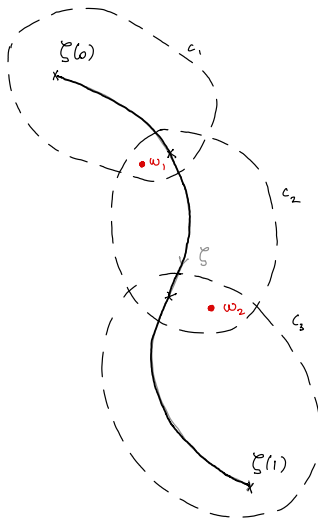
Let $\rho, \iota \in \mathfrak{S}_n$. We define

$$\omega_{\rho, \iota} = (\rho(1) + \mathbf{i}\iota(1), \dots, \rho(n) + \mathbf{i}\iota(n)) \in OC_n$$

$$\omega_{(12)(34), (143)} = \begin{array}{cccc} \cdot & \bullet_1 & \cdot & \cdot \\ \cdot & \cdot & \bullet_4 & \cdot \\ \bullet_2 & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \bullet_3 \end{array}$$

Linearization of ζ

For each c_i, c_{i+1} , we compute ρ, ι such that $\omega_i = \omega_{\rho, \iota}$ lies in the intersection $c_i \cap c_{i+1}$ (Hint: total order extending R and I).



Step 2: linearize ζ

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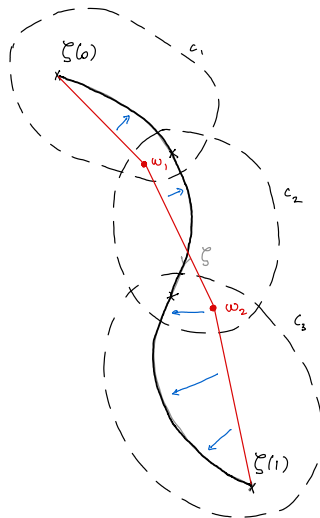
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Linearization of ζ

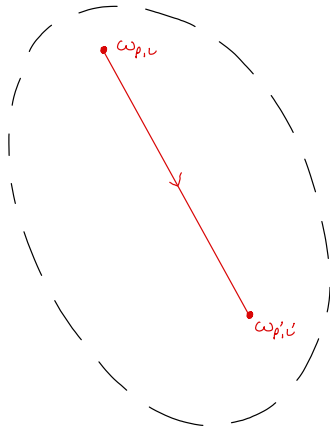
For each c_i, c_{i+1} , we compute ρ, ι such that $\omega_i = \omega_{\rho, \iota}$ lies in the intersection $c_i \cap c_{i+1}$ (Hint: total order extending R and I). The linear interpolation of the ω_i is homotopic to ζ . Why? m.s.c cells are convex!



Step 3: decomposition of the linearization

Reduction

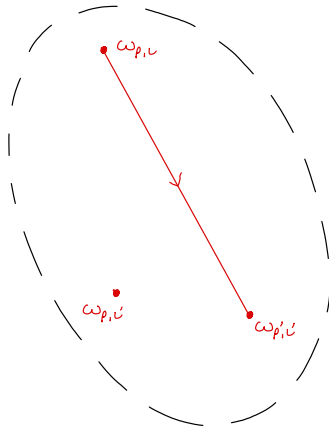
- Computing the braid associated to the whole linearization or to each piece and concatenating the results is equivalent



Step 3: decomposition of the linearization

Reduction

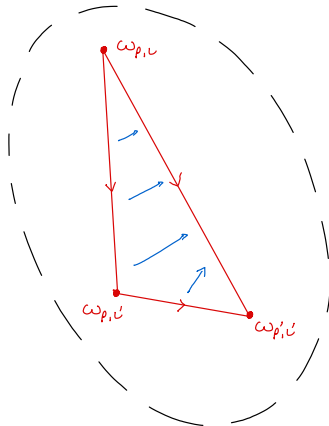
- Computing the braid associated to the whole linearization or to each piece and concatenating the results is equivalent
- Assume $\omega_{\rho,\iota}$ and $\omega_{\rho',\iota'}$ both lie in a m.s.c cell $c = (R, I)$. It means that ρ, ρ' extend R and ι, ι' extend I . **So $\omega_{\rho,\iota'}$ also lies in c !**



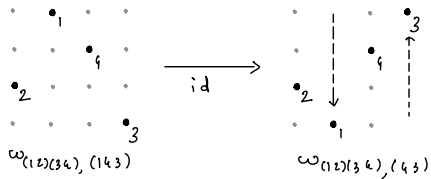
Step 3: decomposition of the linearization

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- We compute the braid of $\omega_{\rho,\iota} \rightarrow \omega_{\rho,\iota'}$ then the braid of $\omega_{\rho,\iota'} \rightarrow \omega_{\rho',\iota'}$



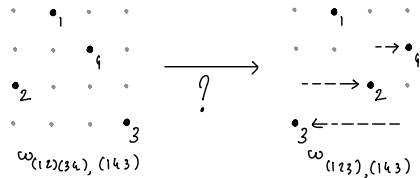
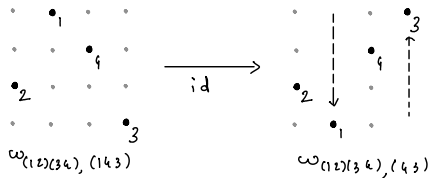
Step 3: decomposition of the linearization



$$\omega_{\rho, \iota} \rightarrow \omega_{\rho, \iota'}$$

The induced braid is trivial, as the real part of the strands is constant.

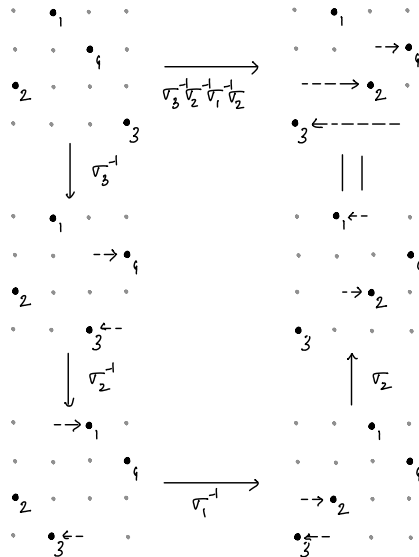
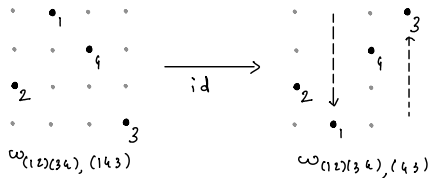
Step 3: decomposition of the linearization



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Step 3: decomposition of the linearization



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The induced braid is trivial, as the real part of the strands is constant.

$$\omega_{\rho, \iota'} \rightarrow \omega_{\rho', \iota'}$$

Let $\rho' \rho^{-1} = s_{i_1} \cdots s_{i_r}$ be a decomposition in elementary transpositions. Output $\sigma_{i_1}^{\varepsilon_1} \cdots \sigma_{i_r}^{\varepsilon_r}$ with $\varepsilon_1, \dots, \varepsilon_r \in \{\pm 1\}$ computed using ι' .

Conclusion

```
~/2025/code/braid_group cargo run --release
```

```
Finished `release` profile [optimized] target(s) in 0.08s
```

```
Running `target/release/braid_group`
```

```
057011025029047051055061063083030504103103500-1010037073097098027049087097-10590150203-1096-1077092
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```

